# CSCI-564 CONSTRAINT PROCESSING AND HEURISTIC SEARCH 

LECTURE 17 - ADVERSARY SEARCH (CONT'D)

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## Recap

- In Game Theory:
- We want to calculate the optimal strategy.
- A strategy (or policy) maps each state to an action.
- Tree Search:
- We limit the search in depth.
- More important to evaluate all the nodes before a certain depth.
- Two algorithms:
- Minimax and negmax.
- Similar ideas.


## $\alpha \beta$-Pruning

- Negmax performs a depth-first search of the game.
- Explore the entire tree until a fixed depth.
- We fall in the same issue we had before.
- The tree can grow very fast.
(6)
-6
(6)
-6 -5 -3 0
(4)

-4
(9)


## $\alpha \beta$-Pruning

- $\alpha \beta$-pruning is a branch-and-bound method.
- The idea is simple:
- Computing the root value while avoiding a part of the tree.
(6)
- Do you have an idea?



## $\alpha \beta$-Pruning

- $\alpha \beta$-pruning uses two bounds called $\alpha \beta$-window:
- $\alpha$ : The least value that the agent can achieve.
- $\beta$ : The maximum value that the opponent can achieve.
- If the initial window is $(-\infty, \infty)$, it will determine the correct value



## $\alpha \beta$-Pruning

- It avoid nodes by achieving cut-offs.
- We can distinguish two types:
- Shallow
- Deep cut-offs


## $\alpha \beta$-Pruning

- Shallow cut-off.
$\square$

You explore the branches. And get the value -8.



## $\alpha \beta$-Pruning

- Shallow cut-off.

The root will be at least equal to 8
a $\square$

You explore the branches. And get the value -8.

c
$\square$
$\square$

## $\alpha \beta$-Pruning

- Shallow cut-off.

The root will be at least equal to 8
a $\square$

You explore the branches. And get the value -8.

$\square$ c


You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Shallow cut-off.

The root will be at least equal to 8
a $\square$

You explore the branches. And get the value -8.
$\square$
b -8
$-5 \quad \mathrm{c} \longleftarrow$ The node will be at least -5 .
$\square$

## $\alpha \beta$-Pruning

- Shallow cut-off.

The root will be at least equal to 8
a $\square$

You explore the branches. And get the value -8.


## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a


## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a


You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a

e

$\geq-5$


You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a



You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a

b
-8

$$
\geq-5
$$


e

```
\geq-5
```

Can already achieve 8 at the root by moving to $b$, so he will never choose to move to e from d.

You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a

b $-8$

$$
\geq-5
$$

$\square$ c


Can already achieve 8 at the root by moving to $b$, so he will never choose to move to e from d.

We can cut the remaining nodes to explore.
You find the value 5 for the node.

## $\alpha \beta$-Pruning

- Deep cut-off.

The root will be at least equal to 8
a


In deep pruning, the bound used for the cut-off can stem not only from the parent, but from any ancestor node.
b $-8$

$$
\geq-5
$$

Can already achieve 8 at the root by moving to b , so he will never choose to move to e from d.

We can cut the remaining nodes to explore.
You find the value 5 for the node.

## $\alpha \beta$-Pruning

## Procedure NegmaxAlphaBeta

Input: Position $u$, bounds $\alpha, \beta$
Output: Value at root
if (leaf(u)) return Eval(u)
$r e s \leftarrow \alpha$
for each $v \in \operatorname{Succ}(u)$
val $\leftarrow-$ NegmaxAlphaBeta( $v,-\beta,-r e s)$
if (val > res) res $\leftarrow$ val
if $(r e s \geq \beta$ ) return res
return res
;; No successor, static evaluation ;; Initialize value res for current frame
;; Traverse successor list
;; Initialize cut-off value ;; Update res
;; Perform cut-off
;; Return final evaluation

## $\alpha \beta$-Pruning

- For any node $u$ :
- $\quad \beta$ is the upper bound used to restrict the node below $u$.
- A cut-off occurs when $u \geq \beta$.
- From the opponent point of view:
- Can choose a move that avoids $u$ with a value no greater then $\beta$.
- The alternate move is no worse than $u$, so searching below $p$ is not necessary.

[^0];; No successor, static evaluation ; Initialize value res for current frame ;; Traverse successor list
;; Initialize cut-off value
;; Update res
;; Perform cut-off
; Return final evaluation

## $\alpha \beta$-Pruning

```
Procedure MinimaxAlphaBeta
Input: Position \(u\), value \(\alpha\), value \(\beta\)
Output: Value at root
if (leaf(u)) return \(\operatorname{Eval}(p)\)
if (max-node(u))
    \(r e s \leftarrow \alpha\)
    for each \(v \in \operatorname{Succ}(u)\)
        val \(\leftarrow \operatorname{MinimaxAlphaBeta(v,res,~} \beta\) )
        \(r e s \leftarrow \max \{r e s, v a l\}\)
        if (res \(\geq \beta\) )
            return res
else
    \(r e s \leftarrow \beta\)
    for each \(v \in \operatorname{Succ}(u)\)
        val \(\leftarrow\) MinimaxAlphaBeta( \(v, \alpha\), res \()\)
        res \(\leftarrow \min \{r e s\), val \(\}\)
        if (res \(\leq \alpha\) )
            return res
return res
```

;; No successor, return evaluation
;; MAX node
;; Initialize result value
;; Traverse successor list
;; Recursion for $\alpha$
;; Take maximal value
;; Result exceeds threshold
;; Propagate value
;; MIN node
;; Initialize result value
;; Traverse successor list
;; Recursion for $\beta$
;; Take minimal value
;; Result exceeds threshold
;; Propagate value
;; Propagate value

## $\alpha \beta$-Pruning



## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
Output: Value at root
if (leaf(u)) return $\operatorname{Eval}(p)$
if (max-node(u))
$r e s \leftarrow \alpha$
for each $v \in \operatorname{SucC}(u)$
val $\leftarrow$ MinimaxAlphaBeta $(v$, res, $\beta$ ) $r e s \leftarrow \max \{r e s, v a /\}$
if $(r e s \geq \beta$ )
return res
else
res $\leftarrow \beta$
for each $v \in \operatorname{Succ}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, $\alpha$, res $)$
res $\leftarrow \min \{r e s, v a l\}$
if $(r e s \leq \alpha)$
return res
return res
;; No successor, return evaluation
; MAX node
;; Initialize result value
; Traverse successor list
;: Recursion for $\alpha$
;; Take maximal value ; Result exceeds threshold ;; Propagate value ; MIN node ; Initialize result value ; Traverse successor list ;; Recursion for $\beta$
;; Take minimal value ; Result exceeds threshold ;; Propagate value ;; Propagate value


## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
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if $(r e s \geq \beta$ )
return res

## else

res $\leftarrow \beta$
for each $v \in \operatorname{Succ}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, $\alpha$, res) $r e s \leftarrow \min \{r e s, v a l\}$
if $(r e s \leq \alpha)$
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;; Propagate value
; MIN node
; Initialize result value
; Traverse successor list ;; Recursion for $\beta$
;: Take minimal value ;; Result exceeds threshold
;; Propagate value
; Propagate value

## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
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if $(r e s \geq \beta$ )
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if $(r e s \leq \alpha)$
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; Take minimal value ; Result exceeds threshold
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; Propagate value

## 4

## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
Output: Value at root
if (leaf(u)) return Eval(p)
if (max-node(u))

```
res \(\leftarrow \alpha\)
    for each \(v \in \operatorname{Succ}(u)\)
        val \(\leftarrow\) MinimaxAlphaBeta \((v, r e s, \beta)\)
        \(r e s \leftarrow \max \{r e s, v a l\}\)
            if \((r e s \geq \beta)\)
            return res
else
    \(r e s \leftarrow \beta\)
    for each \(v \in \operatorname{Succ}(u)\)
        val \(\leftarrow\) MinimaxAlphaBeta(v, \(\alpha\), res)
        \(r e s \leftarrow \min \{r e s, v a l\}\)
        if \((r e s \leq \alpha)\)
        return res
return res
return res
```

;; No successor, return evaluation
;; MAX node
;; Initialize result value
;; Traverse successor list
i; Recursion for $\alpha$
;; Take maximal value ; Result exceeds threshold ;; Result exceeds threshold
;; Propagate value ; MIN node ;; Initialize result value ; Traverse successor list ;; Recursion for $\beta$ ;; Take minimal value

4

## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
Output: Value at root
if (leaf(u)) return Eval(p)
if (max-node(u))
res $\leftarrow \alpha$
for each $v \in \operatorname{SucC}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, res, $\beta$ ) $r e s \leftarrow \max \{r e s, v a /\}$
if $(r e s \geq \beta$ )
return res

## else

res $\leftarrow \beta$
for each $v \in \operatorname{Succ}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, $\alpha$, res) $r e s \leftarrow \min \{r e s, v a l\}$
if $(r e s \leq \alpha)$
return res
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;; No successor, return evaluation
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## $\alpha \beta$-Pruning

| Procedure MinimaxAlphaBeta Input: Position $u$, value $\alpha$, value $\beta$ Output: Value at root |  |
| :---: | :---: |
| if (leaf(u)) return $\operatorname{Eval}(p)$ | ;; No successor, return evaluation |
| if (max-node(u)) | ; MAX node |
| $r e s \leftarrow \alpha$ | ;; Initialize result value |
| for each $v \in \operatorname{Succ}(u)$ | ;; Traverse successor list |
| val $\leftarrow$ MinimaxAlphaBeta( $v$, res, $\beta$ ) | ;; Recursion for $\alpha$ |
| $r e s \leftarrow \max \{r e s, v a l\}$ | ;; Take maximal value |
| if (res $\geq \beta$ ) | ;; Result exceeds threshold |
| return res | ;; Propagate value |
| else | ;; MIN node |
| $r e s \leftarrow \beta$ | ;; Initialize result value |
| for each $v \in \operatorname{Succ}(u)$ | ;; Traverse successor list |
| val $\leftarrow$ MinimaxAlphaBeta(v, $\alpha$, res) | ;; Recursion for $\beta$ |
| $r e s \leftarrow \min \{r e s, v a l\}$ | ;; Take minimal value |
| if $(r e s \leq \alpha)$ return res | ;; Result exceeds threshold <br> ;; Propagate value |
| return res | ;; Propagate value |

$$
\leq-1
$$

## 4

e ;; Initialize result value ;; Traverse successor list Take minimal value
value

## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
Output: Value at root
if (leaf(u)) return $\operatorname{Eval}(p)$
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$r e s \leftarrow \alpha$
for each $v \in \operatorname{SucC}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, res, $\beta$ ) $r e s \leftarrow \max \{r e s$, val $\}$
if $(r e s \geq \beta$ )
return res

## else

res $\leftarrow \beta$
for each $v \in \operatorname{Succ}(u)$
val $\leftarrow$ MinimaxAlphaBeta(v, $\alpha$, res) $r e s \leftarrow \min \{r e s, v a l\}$
if $(r e s \leq \alpha)$
return res
return res
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; Recursion for $\alpha$
;; Take maximal value ; Result exceeds threshold
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4
B



## $\alpha \beta$-Pruning

## Procedure MinimaxAlphaBeta

Input: Position $u$, value $\alpha$, value $\beta$
Output: Value at root
if (leaf(u)) return $\operatorname{Eval}(p)$
if (max-node(u))
$r e s \leftarrow \alpha$
for each $v \in \operatorname{Succ}(u)$
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$r e s \leftarrow \max \{r e s, v a /\}$
if $(r e s \geq \beta$ )
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res $\leftarrow \beta$
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$r e s \leftarrow \min \{r e s, v a l\}$
if (res $\leq \alpha$ )
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4
$\leq-1$


$$
\leq-7
$$



## Performance $\alpha \beta$-Pruning

- To obtain an upper bound on the performance we need to prove that there is a minimal tree that must be searched.
- This tree is called critical tree.
- And the nodes are critical nodes.
- The number of leaves on the critical tree is:

$$
b^{\lceil d / 2\rceil}+b^{\lfloor d / 2\rfloor}-1
$$

- Where $b$ is the number of children of a node and $d$ is the height.


## Performance $\alpha \beta$-Pruning

- The number of leaves on the critical tree is:

$$
b^{\lceil d / 2\rceil}+b^{\lfloor d / 2\rfloor}-1
$$

- Where $b$ is the number of children of a node and $d$ is the height.
- Motivation:
- To prove that the value of the root is at least $v, b^{\lceil d / 2\rceil}$ leaves must be inspected.
- One move must be considered on each agent level
- And all moves on each opponent level.
- $b^{\lfloor d / 2\rfloor}-1$ leaf nodes must be generated to see that the value is at most $v$.
- Opponent strategy.


## Performance $\alpha \beta$-Pruning

- The number of leaves on the critical tree is:

$$
b^{\lceil d / 2\rceil}+b^{\lfloor d / 2\rfloor}-1
$$

- Where $b$ is the number of children of a node and $d$ is the height.
- Any remarks?
- $b^{[d / 2]}+b^{\lfloor d / 2]}-1 \approx 2 \sqrt{b^{d}}$
- Which is the number of leaves in the whole unpruned tree.
- Consequently $\alpha \beta$-pruning could double the search.
- If the tree is searched in the best-first order.
- The best move is always chosen first for exploration.
- Then $\alpha \beta$ will search only the critical tree.


[^0]:    Procedure NegmaxAlphaBeta
    Input: Position $u$, bounds $\alpha, \beta$ Output: Value at root
    if (leaf(u)) return Eval(u)
    res $\leftarrow \alpha$
    for each $v \in \operatorname{Succ}(u)$
    val $\leftarrow-$ NegmaxAlphaBeta( $v,-\beta$,-res)
    if (val > res) res $\leftarrow \mathrm{val}$
    if $(r e s \geq \beta)$ return res

